

Math 410 Intro to Quantum Computing Homework 8

Sample Solution

6.1 Let $N = 2^n$, $\tilde{f} = (\tilde{f}(0), \dots, \tilde{f}(N-1))^t$ and $f = (f(0), \dots, f(N-1))^t$. Then

$$\sum_y |\tilde{f}(y)|^2 = \langle \tilde{f} | \tilde{f} \rangle = \langle \tilde{f}(y) | \underbrace{K^\dagger(x, y) K(x, y)}_{\mathbb{I}_N} | \tilde{f}(y) \rangle = \langle f | f \rangle = \sum_x |f(x)|^2 \text{ as } K \text{ is unitary.}$$

6.2 (1) To get $\langle x | x \rangle = 1$, note that $\langle \psi | \psi \rangle = \mathcal{N}^2 \sum_x \cos^2 \frac{2\pi x}{N}$. Thus, $\sum_x \cos^2 \frac{2\pi x}{N} = \sum_x \frac{\cos(\frac{4\pi x}{N}) + 1}{2} = \frac{N}{2}$, and hence, $\mathcal{N} = 2^{-\frac{n-1}{2}}$

(2) The coefficient of $|x\rangle$ in $U_{QFT_n} |\psi\rangle$ equals $\frac{1}{\sqrt{N}\sqrt{N/2}} \sum_y e^{-2\pi i xy/N} \cos \frac{2\pi y}{N}$
 $= \frac{1}{\sqrt{2N}} e^{-2\pi i xy/N} (e^{2\pi i y/N} + e^{-2\pi i y/N}) = \frac{1}{\sqrt{2N}} (e^{-2\pi i(x+1)y/N} + e^{-2\pi i(x-1)y/N})$,
 which is nonzero if and only if $x = 1$ or $N - 1$. So, the expression reduces to $\frac{1}{\sqrt{2N}} N(\delta_{x,1} + \delta_{x,N-1}) = \frac{1}{\sqrt{2}} (\delta_{x,1} + \delta_{x,N-1})$. Therefore, $U_{QFT_n} |\psi\rangle = \frac{1}{\sqrt{2}} (|1\rangle + |N-1\rangle)$.

6.3 Let $\frac{2^n}{P} = m \in \mathbb{N}$. Apply QFT, we obtain $|\Psi'\rangle = \frac{1}{2^n} \sum_{x,y} e^{-2\pi i xy/2^n} |y\rangle |f(x)\rangle$.

Separating the summation over x using an arbitrary function $h(x)$, we let

$$\sum_{x=0}^{2^n-1} h(x) = \sum_{l=0}^{P-1} \sum_{k=0}^{m-1} h(kP+l)$$

Substituting, we have $|\Psi'\rangle = \frac{1}{2^n} \sum_{l=0}^{P-1} \sum_{k=0}^{m-1} \sum_y e^{-2\pi i ly/2^n} e^{-2\pi i ky/m} |y\rangle |f(kP+l)\rangle$

As f has period P , $f(kP+l) = f(l)$, $|\Psi'\rangle = \frac{1}{2^n} \sum_y \sum_{k=0}^{m-1} e^{-2\pi i ky/m} \sum_{l=0}^{P-1} e^{-2\pi i ly/2^n} |y\rangle |f(l)\rangle$.

When $y = qm, 0 \leq q < P-1$, then $\sum_{k=0}^{m-1} e^{-2\pi i ky/m} = \sum_{k=0}^{m-1} e^{-2\pi i kq} = m$;
 When $y \neq qm$, then $\sum_{k=0}^{m-1} e^{-2\pi i ky/m} = \frac{1 - e^{-2\pi i y}}{1 - e^{-2\pi i y/m}} = 0$.
 $1 + a + \dots + a^k = \frac{1 - a^{k+1}}{1 - a}$

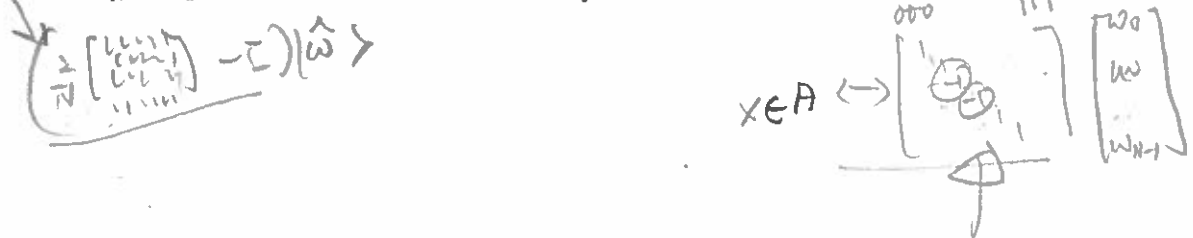
With that, we have $|\Psi'\rangle = \frac{m}{2^n} \sum_{l=0}^{P-1} \sum_{q=0}^{P-1} e^{-2\pi i lq/P} |qm\rangle |f(l)\rangle$.

Therefore, we can obtain $qm = \frac{q(2^n)}{P}$ for some $q \in \mathbb{Z}, 0 \leq q \leq P-1$.

7.1 Similar to Proposition 7.1, we have

$$U_f |\phi\rangle = DR_f |\phi\rangle = D(\sum_{x \neq z} \omega_x |x\rangle - \sum_{z \in A} \omega_z |z\rangle) = \sum_{x \notin A} (2\bar{\omega} - \omega_x) |x\rangle + \sum_{z \in A} (2\bar{\omega} + \omega_z) |z\rangle$$

Alternatively, one can look at the vectors with components $|w\rangle = (w_0, \dots, w_{N-1})^t$ and see that $|\bar{w}\rangle = R_f(|w\rangle)$ will change the entries of $|w\rangle$ with labels in A to negative, and then $D|\bar{w}\rangle = \frac{2}{N} \bar{w} \sum_x |x\rangle - \bar{w}$. Thus, the entries of $U_f |w\rangle$ has the asserted form.



7.2 By induction on $k \geq 0$. When $k = 0$, the result holds. Assume that the result holds for $k - 1$. Then $U_f^k |\varphi_0\rangle = U_f |\varphi_{k-1}\rangle$. By induction assumption $|\varphi_{k-1}\rangle = (w_0, \dots, w_{N-1})^t$, where $w_x = a_{k-1}$ if $x \in A$ and $w_x = b_{k-1}$ if $x \notin A$. Hence, $|\varphi_k\rangle = U_f(|\varphi_{k-1}\rangle)$ has two types of entries depending on whether the entries have label in A or not:

$$a_k = \frac{2}{N}((N-d)b_{k-1} - da_{k-1}) + a_{k-1} = \frac{1}{N}((N-2d)a_{k-1} + 2(N-d)b_{k-1}),$$

$$b_k = \frac{2}{N}((N-d)b_{k-1} - da_{k-1}) - b_{k-1} = \frac{1}{N}(-2da_{k-1} + (N-2d)b_{k-1}).$$

Remark Note the typo in the book. We should have $a_k = \frac{1}{N}((N-2d)a_{k-1} - 2(N-d)b_{k-1})$.

7.3 Similar to proposition 7.3, we have

Let $p_k = \sqrt{d}a_k, q_k = \sqrt{N-d}b_k$. Then $\begin{bmatrix} p_k \\ q_k \end{bmatrix} = M \begin{bmatrix} p_{k-1} \\ q_{k-1} \end{bmatrix}$, where

$$M = \begin{bmatrix} (N-2d)/N & 2\sqrt{N-d}/N \\ -2\sqrt{N-d}/N & (N-2d)/N \end{bmatrix} = \begin{bmatrix} \cos 2\theta & \sin 2\theta \\ -\sin 2\theta & \cos 2\theta \end{bmatrix} \rightarrow \begin{matrix} \text{cos & sin} \\ \text{-sin & cos} \end{matrix}$$

Then $\begin{bmatrix} p_k \\ q_k \end{bmatrix} = M^k \begin{bmatrix} p_0 \\ q_0 \end{bmatrix} = \begin{bmatrix} \cos 2k\theta & \sin 2k\theta \\ -\sin 2k\theta & \cos 2k\theta \end{bmatrix} \begin{bmatrix} \sin \theta \\ \cos \theta \end{bmatrix} = \begin{bmatrix} \sin[(2k+1)\theta] \\ \cos[(2k+1)\theta] \end{bmatrix}$

As $a_k = \frac{1}{\sqrt{d}}p_k, b_k = \frac{1}{\sqrt{N-d}}q_k$ we have $a_k = \frac{1}{\sqrt{d}} \sin[(2k+1)\theta], b_k = \frac{1}{\sqrt{N-d}} \cos[(2k+1)\theta]$.

7.4 Similar to proposition 7.4, we have

Define \tilde{m} by $(2\tilde{m}+1)\theta = \frac{\pi}{2} \rightarrow \tilde{m} = \frac{\pi}{4\theta} - \frac{1}{2}$. We see $|\tilde{m} - m| \leq \frac{1}{2}$ as $m = \lfloor \frac{\pi}{4\theta} \rfloor$.

Then $|(2m+1)\theta - (2\tilde{m}+1)\theta| = |(2m+1)\theta - \frac{\pi}{2}| \leq \theta$.

Note that $\theta \sim \frac{d}{\sqrt{N}}$ is small when $N \gg 1$ and $\sin x$ is monotonically increasing in neighborhood of $x = 0$. Thus, $0 < \sin |(2m+1)\theta - \frac{\pi}{2}| < \sin \theta$ or $\cos^2 |(2m+1)\theta| \leq \sin^2 \theta = \frac{d}{N}$.

Hence, $P_{m,z} = \sin^2[(2m+1)\theta] = 1 - \cos^2[(2m+1)\theta] \geq 1 - \frac{d}{N}$ asserted in (7.47)

Eventually, $m = \lfloor \frac{\pi}{4\theta} \rfloor \leq \frac{\pi}{4\theta} \leq \frac{\pi}{4} \sqrt{\frac{N}{d}}$; so, the operation time of m is $O(\sqrt{\frac{N}{d}})$.

Check: $\begin{bmatrix} a_0 \\ b_0 \end{bmatrix} = \begin{bmatrix} 1 \\ 1 \end{bmatrix} \frac{1}{\sqrt{N}}$
 $\begin{bmatrix} p_0 \\ q_0 \end{bmatrix} = \begin{bmatrix} 1 \\ 1 \end{bmatrix} \frac{1}{\sqrt{N}}$
 with yield the correct correspondence $\begin{bmatrix} p_1 \\ q_1 \end{bmatrix}$
 $\begin{bmatrix} p_1 \\ q_1 \end{bmatrix}$

$$\rho = \begin{bmatrix} a & \sqrt{a-b} \\ \sqrt{a-b} & b \end{bmatrix} = \frac{\begin{bmatrix} 1 & 1 \\ 1 & 1 \end{bmatrix}}{2}$$

$$\rho_S = \begin{pmatrix} a & b \\ b & c \end{pmatrix}$$

$$I_2 = \frac{\begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix}}{2}$$

$$L(\rho) = \left(\frac{1-p}{2}\right) I_2$$

Depolarizing Channel

The depolarizing channel is defined by

$$\mathcal{E}(\rho_S) = (1-p)\rho_S + pI/2, \quad 0 \leq p \leq 1.$$

The input state ρ_S is sent to a maximally mixed state $I/2$ with a probability p , remains unchanged with a probability $1-p$. The Choi/Kraus operators are: $E_0 = \sqrt{1-3p/4}I$ and $E_k = \sqrt{p/4}\sigma_k$ for $k = x, y, z$. Let $p' = 3p/4$. Then

$$\mathcal{E}(\rho_S) = (1-p')\rho_S + (p'/3) \sum_k \sigma_k \rho_S \sigma_k$$

$$= (\sqrt{1-p'} \cdot I) \rho_S (\sqrt{1-p'} \cdot I)$$

$$\sigma_x = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}$$

$$\sigma_y = \begin{bmatrix} 0 & -i \\ i & 0 \end{bmatrix}$$

$$\sigma_z = \begin{bmatrix} 1 & 0 \\ 0 & -1 \end{bmatrix}$$

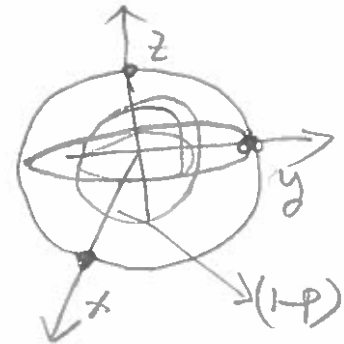
This can be modeled by

$$\rho_S \mapsto V(\rho_S \otimes I/2 \otimes [(1-p)|0\rangle\langle 0| + p|1\rangle\langle 1|])V^\dagger = (1-p)\rho_S \otimes (I/2) \otimes |0\rangle\langle 0| + p(I/2) \otimes \rho_S \otimes |1\rangle\langle 1|$$

with $V \equiv I_4 \otimes |0\rangle\langle 0| + U_{SWAP}|1\rangle\langle 1|$.

Use the Bloch sphere representation:

$$\mathcal{E}(\rho_S) = \frac{1}{2} \left(I + (1-p) \sum_k c_k \sigma_k \right)$$



So, the Bloch sphere shrinks in the x, y, z directions by a factor of $1-p$.

$$\mathcal{E} \left(\frac{I}{2} + c_x \sigma_x + c_y \sigma_y + c_z \sigma_z \right)$$

$$= \left(\frac{I}{2} + (1-p)(c_x \sigma_x + c_y \sigma_y + c_z \sigma_z) \right)$$

$$\mathcal{E} \left(\frac{I}{2} + \frac{\sigma_x}{2} \right)$$

$$\mathcal{E} \left(\frac{I}{2} + \frac{\sigma_y}{2} \right)$$

$$\mathcal{E} \left(\frac{I}{2} + \frac{\sigma_z}{2} \right)$$

$$= \begin{pmatrix} 1 & 0 \\ 0 & 1-p \end{pmatrix} + \begin{pmatrix} 0 & 0 \\ 0 & p \end{pmatrix}$$

$$\rightarrow E_0^\dagger E_0 + E_1^\dagger E_1 = I_2$$

Amplitude-Damping Channel

The amplitude-damping channel is defined by

$$\mathcal{E}(\rho_S) = E_0 \rho_S E_0^\dagger + E_1 \rho_S E_1^\dagger, \quad E_0 = \begin{pmatrix} 1 & 0 \\ 0 & \sqrt{1-p} \end{pmatrix}, \quad E_1 = \sqrt{p} \begin{pmatrix} 0 & 1 \\ 0 & 0 \end{pmatrix}$$

It describes the decay of the qubit from $|1\rangle$ to $|0\rangle$ with probability p .

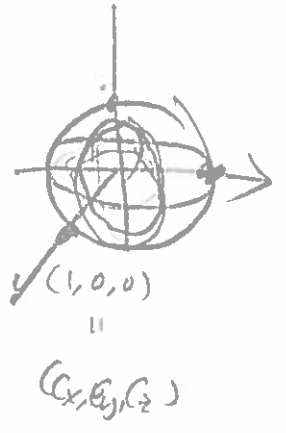
Use the block sphere representation:

$$\mathcal{E}(\rho_S) = \frac{1}{2} (I + \sqrt{1-p} c_x \sigma_x + \sqrt{1-p} c_y \sigma_y + [p + (1-p)c_z] \sigma_z)$$

So, the block sphere shrinks in the x and y directions by a factor of $\sqrt{1-p}$, and shrinks in the z direction by a factor of $1-p$ and then shift to the north pole by p .

$$\rho_S = \frac{1}{2} (I + c_x \sigma_x + c_y \sigma_y + c_z \sigma_z)$$

$$\mathcal{E}(\rho_S) = \frac{1}{2} (I + \sqrt{1-p} c_x \sigma_x + \sqrt{1-p} c_y \sigma_y + [p + (1-p)c_z] \sigma_z)$$



9.4 Lindblad Equation

Recall the Liouville-von Neumann equation for a closed system

$$\frac{\partial \rho_S}{\partial t} = \frac{1}{i} [H, \rho_S].$$

For an open system, it becomes

$$\rho_S(t) = \text{Tr}_E(U(t)(\rho_S \otimes \rho_E)U(t)^\dagger).$$

Thus,

$$\frac{\partial \rho_S}{\partial t} = \frac{\partial \rho_S}{\partial t} \text{Tr}_E(U(t)(\rho_S \otimes \rho_E)U(t)^\dagger) = \frac{1}{i} \text{Tr}_E[H, U(t)(\rho_S \otimes \rho_E)U(t)^\dagger].$$

Assume that the system relies very little about previous history, then we use the **Markovian approximation**, and the equation will be of the form

$$\boxed{\frac{\partial \rho_S}{\partial t}} = \mathcal{L}\rho_S(t) = \frac{1}{i} [\tilde{H}, \rho_S] + \sum_{K=1}^N \gamma_K \left(L_K \rho_S L_K^\dagger - \frac{1}{2} L_K^\dagger L_K \rho_S - \frac{1}{2} \rho_S L_K^\dagger L_K \right).$$

This is known as the Lindblad equation; the operators L_k are known as the Lindblad operators.

Quantum error correction

10.1 Classical quantum error correction

In classical communication, suppose a noisy channel will change x to $x \oplus 1$ for $x = 0, 1$ with a probability $p < 1/2$. One may send xxx instead of x and use maximum likelihood decoding so that the probability of correct transmission for one bit is the sum of the probabilities of x sent to xxx , $x\bar{x}$, $\bar{x}x$, $\bar{x}\bar{x}$:

~~xxx~~ $(1-p)^3 + 3p(1-p)^2 = (1-p)^2(1+2p) \gg p$

So if Alice send xxx , then Bob decodes the message as xxx using the above scheme. i.e. $\frac{(1-p)^3 + 3p(1-p)^2}{(1-p)^2(1+2p)}$

Without encoding, Bob will correctly decode with a prob. $1-p$

0, 1

$C = \{(000), (111)\}$
 $\subseteq \mathbb{Z}_2^3$ with 8 elements

$$(1-p)(1+2p) = 1+p-2p^2 > 1$$

$$(1-p)^2(1+2p) > (1-p)$$

10.2 Quantum error correction

10.2.1 Bit-Flip QECC

We cannot copy qubit because of no-cloning.

But, we can encode $|\psi\rangle = a|0\rangle + b|1\rangle$ as $|\psi\rangle_L = a|000\rangle + b|111\rangle$ using *CNOT* gates on $|\psi\rangle|00\rangle$.

Then the set of **code words** is

$$C = \{a|000\rangle + b|111\rangle : a, b \in \mathbb{C}, |a|^2 + |b|^2 = 1\}$$

The QECC scheme can be done in the following steps.

- Encode $|\psi\rangle = a|0\rangle + b|1\rangle$ as $|\psi\rangle_L = a|000\rangle + b|111\rangle \in C$.

- Transmit via the quantum channel.

- Apply error syndrome detection and correction

* Suppose $|x_1 x_2 x_3\rangle$ is received.

* Add two ancillas $|AB\rangle$ with $A = |x_1 \oplus x_2\rangle$ and $B = |x_1 \oplus x_3\rangle$ to detect the (error) syndrome.

* Apply correction to correct $|\psi\rangle_L$ accordingly.

- Reversing the encoding step, one gets $|\psi\rangle$.

$|\psi\rangle = a|0\rangle + b|1\rangle$
 "copy" means $|\psi\rangle|00\rangle = |a\rangle|b\rangle|a\rangle|b\rangle$

$$E(|\psi\rangle) = (P_0^x)^p + b(X^0 I^0 I^0)P_0^x + (I^0 X^1 I^0)P_0^x + (I^0 I^0 X^1)P_0^x + P_1^x$$

Bit-flip channel.

$$\begin{pmatrix} a \\ b \end{pmatrix} \mapsto X \begin{pmatrix} a \\ b \end{pmatrix} = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} \begin{pmatrix} a \\ b \end{pmatrix} = \begin{pmatrix} b \\ a \end{pmatrix}$$