

# Amidakuji, sorting algorithms, & permutations:

From recreational mathematics to research mathematics

Chi-Kwong Li (ckli@math.wm.edu)  
Department of Mathematics  
The College of William and Mary

# Amidakuji/Ghost Leg Drawing

## Amidakuji

At first, you see a group of lines at the top and the same number of lines at the bottom.

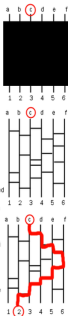
The middle is covered up so you can't tell which top line leads to which bottom line.

1. Everyone chooses or is assigned a top line.

2. The bottom lines are assigned to things to be distributed, such as prizes or duties.

3. The amidakuji is revealed.

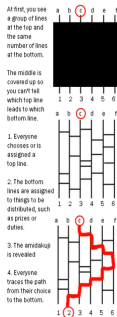
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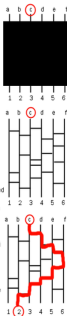
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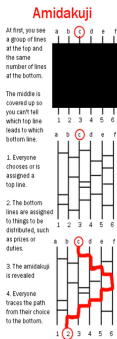
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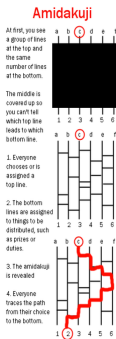
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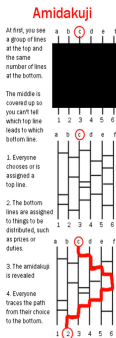
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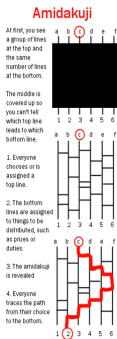
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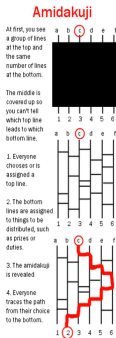
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## Questions

- Why do we always get an one-one correspondence (bijection)?
- Can we get all possible job assignments?
- What is the minimum number of horizontal segments needed for a given job assignment?

# Answer of Question 1

## **George Polya (1887-1985)**

If one cannot solve a problem,  
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- What if there is no horizontal line segment?
- What if there is one horizontal line segment?
- An easy induction argument!

- Record the job assignment as a permutation (a seat reassignment), say,

$$\sigma = [5, 3, 1, 2, 4] = \begin{pmatrix} 1 & 2 & 3 & 4 & 5 \\ 5 & 3 & 1 & 2 & 4 \end{pmatrix}.$$

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- Use Coxeter transpositions  $(1, 2), (2, 3), (3, 4), (4, 5)$  to restore order.
- Key idea** Determine the number  $\iota(\sigma)$  of **inversions** of  $\sigma$ .



**Example** For  $\sigma = [5, 3, 1, 2, 4]$ , total number of inversions is:

$$\iota(5) + \iota(4) + \iota(3) + \iota(2) + \iota(1) = 4 + 0 + 2 + 0 + 0 = 6,$$

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$$(n-1) + \dots + 1 = n(n-1)/2 \text{ steps.}$$

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**Example.**  $\sigma = \begin{pmatrix} 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 \\ 3 & 4 & 5 & 9 & 6 & 7 & 1 & 8 & 2 \end{pmatrix} = (1, 3, 5, 6, 7)(2, 4, 9)(8).$

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Then  $\sigma = (1, 7)(1, 6)(1, 5)(1, 3)(2, 9)(2, 4).$

# Some open problems

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- One can write a computer program to determine  $r^*$ .

# Partial results of the general problem

We have the following list for  $r^*(n, m)$  for  $S_n$  and  $(i, i + \ell)$  with  $\ell \leq m$ ,

$n \setminus m$	1	2	3	4	5	6	7	8	9	10	11
2	1										
3	3	2									
4	6	4	3								
5	10	5	5	4							
6	15	[7]	6	6	5						
7	21	[10]	8	7	7	6					
8	28	[14]	[10]	9	8	8	7				
9	36	[16]	[11]	10	10	9	9	8			
10	45	[19]	[14]	[12]	11	11	10	10	9		
11	55	[23]	[16]	[14]	13	12	12	11	11	10	
12	66	29*	20*	17*	16*	14	13	13	12	12	11

where the entries marked by brackets are obtained by computer programming.

- Theoretical computer science.

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# Related research

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Study the change of genetic sequences  $x_1x_2x_3 \cdots$ , with  $x_i \in \{A, U, G, C\}$ .

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- **Quantum computing.**

Decompose certain quantum gates into simple quantum gates (CNOT gates).



Reference:

Zejun Huang, Chi-Kwong Li, Sharon H. Li, Nung-Sing Sze, Factorization of permutations, preprint, arXiv:1303.3776.

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**Thank you for your attention!**

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**Open problem** What is the formula?

- If we use  $(1, 2)$ , the forward and backward long cycles  $(1, \dots, n)$  and  $(1, n, n-1, \dots, 2)$ , then  $r^*$  is given by:

$S_2$	$S_3$	$S_4$	$S_5$	$S_6$	$S_7$	$S_8$	$S_9$	$S_{10}$	$S_{11}$
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**Conjecture**  $r^* = \binom{n}{2}$  for  $n \geq 4$ .