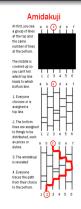
# Amidakuji (ghost leg): From a simple game to research in different areas

Chi-Kwong Li
The College of William and Mary



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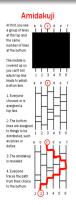
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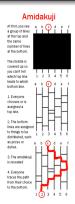
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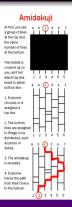
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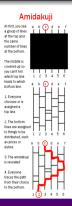


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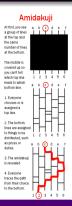


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- Why do we always get an one-one correspondence (bijection)?
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- What is the minimum number of horizontal segments needed for a given job assignment?

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- An easy induction argument!

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$$(n-1) + \cdots + 1 = n(n-1)/2$$
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**Example.** 
$$\sigma = \begin{pmatrix} 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 \\ 3 & 4 & 5 & 9 & 6 & 7 & 1 & 8 & 2 \end{pmatrix} = (1,3,5,6,7)(2,4,9)(8).$$



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Then  $\sigma = (1,7)(1,6)(1,5)(1,3)(2,9)(2,4)$ .



## Open problems

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- Determine the optimal (smallest)  $r^* = r^*(n, m)$  so that every  $\sigma \in S_n$  is a product at most  $r^*$  transpositions in  $G_m$ .

## Partial results of the general problem

We have the following list for  $r^*(n,m)$  for  $S_n$  and  $(i,i+\ell)$  with  $\ell \leq m$ ,

$n \backslash m$	1	2	3	4	5	6	7	8	9	10
2	1									
3	3	2								
4	6	4	3							
5	10	5	5	4						
6	15	[7]	6	6	5					
7	21	[10]	8	7	7	6				
8	28	[14]	[10]	9	8	8	7			
9	36	[16]	[11]	10	10	9	9	8		
10	45	[19]	[14]	[12]	11	11	10	10	9	
11	55	24?	18?	15?	13	12	12	11	11	10

where the entries marked by brackets are obtained by computer programming.

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$S_2$	$S_3$	$S_4$	$S_5$	$S_6$	$S_7$	$S_8$	$S_9$	$S_{10}$	$S_{11}$	$S_{12}$
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• Diameter of the Cayley graph of  $S_n$  using L, S, and  $R=L^{-1}=(n,n-1,\ldots,1)$ :

	$S_2$	$S_3$	$S_4$	$S_5$	$S_6$	$S_7$	$S_8$	$S_9$	$S_{10}$	$S_{11}$
ĺ	1	2	6	10	15	21	28	36	45	???

#### Computing

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$$\begin{pmatrix} I_{j-1} & & & \\ & Q & & \\ & & I_{n-j-1} \end{pmatrix}, \quad \text{ with } Q \in M_2, \qquad j=1,\dots,n-1.$$

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For n=4, we need

$$\begin{pmatrix} * & * & 0 & 0 \\ * & * & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix}, \quad \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & * & * & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix}, \quad \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & * & * \\ 0 & 0 & * & * \end{pmatrix}.$$

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- More on Quantum computing.
   It is of interest to decompose certain quantum gates into simpler quantum gates (CNOT gates).

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Thank you for your attention!