

# Math 410 Intro to Quantum Computing Homework 10

## Sample Solution

10.2 (a)  $U_H Z U_H = \left(\frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix}\right) \begin{bmatrix} 1 & 0 \\ 0 & -1 \end{bmatrix} \left(\frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix}\right) = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} = X.$

(b)  $\begin{bmatrix} 1 & 0 \\ 0 & -1 \end{bmatrix} \frac{1}{\sqrt{2}}(|0\rangle + |1\rangle) = \frac{1}{\sqrt{2}}(|0\rangle - |1\rangle) = |-\rangle$ ;  $\begin{bmatrix} 1 & 0 \\ 0 & -1 \end{bmatrix} \frac{1}{\sqrt{2}}(|0\rangle - |1\rangle) = \frac{1}{\sqrt{2}}(|0\rangle - (-|1\rangle)) = |+\rangle$ ;

(c)  $\begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} \frac{1}{\sqrt{2}}(|0\rangle + |1\rangle) = \frac{1}{\sqrt{2}}(|0\rangle + |1\rangle) = |+\rangle$ ;  $\begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} \frac{1}{\sqrt{2}}(|0\rangle - |1\rangle) = \frac{1}{\sqrt{2}}(-|0\rangle + |1\rangle) = -|-\rangle$ ;

10.3 Note that  $e^{i\beta Z} = \cos \beta I + i \sin \beta Z$ , which is a linear combination of  $I$  and  $Z$ . So, the linear span of  $E = \{IIZ, IZI, IIZ\}$  is the same as that of  $E' = \{IIU_\beta, IU_\beta IU_\beta II\}$ . Hence, the QECC for channel with error operators in  $E$  also works for error operators in  $E'$ .

Alternatively, suppose  $U_\beta$  occurs in the first qubit. Then  $(U_\beta \otimes I \otimes I)(a|+++ \rangle + b|--- \rangle)$   
 $= a(\cos \beta|+++ \rangle + i \sin \beta| - ++ \rangle) + b(\cos \beta| --- \rangle + \sin \beta| + -- \rangle).$

The Hadamard gates in the error syndrome detection circuit maps these vectors as

$$a(\cos \beta|000\rangle + i \sin \beta|100\rangle) + b(\cos \beta|111\rangle + i \sin \beta|011\rangle).$$

The error syndrome detection circuit outputs  $\cos \beta(a|000\rangle + b|111\rangle)|00\rangle + i \sin \beta(a|100\rangle + b|011\rangle)|11\rangle$ . The receiver gets  $a|000\rangle + b|111\rangle$  when they measure ancilla qubits 00 and 11 with probabilities  $\cos^2 \beta$  and  $\sin^2 \beta$ , respectively. So, the circuit in Figure 10.4 can correct the error, and the Hadamard gates will give the correct transmitted state.

10.4 If  $|\psi\rangle = |0\rangle$ , there is  $\frac{1}{\sqrt{2}}$  chance after the Hadamard gate it is mapped into  $|0\rangle$ , and after triparting the system becomes  $|000\rangle$ ; there is  $\frac{1}{\sqrt{2}}$  chance after the Hadamard gate it is mapped into  $|1\rangle$ , and after triparting the system becomes  $|111\rangle$ . Therefore,  $|0\rangle|00\rangle \rightarrow \frac{1}{\sqrt{2}}(|000\rangle + |111\rangle) = |+\rangle$

If  $|\psi\rangle = |1\rangle$ , there is  $\frac{1}{\sqrt{2}}$  chance after the Hadamard gate it is mapped into  $|0\rangle$ , and after triparting the system becomes  $|000\rangle$ ; there is  $\frac{1}{\sqrt{2}}$  chance after the Hadamard gate it is mapped into  $-|1\rangle$ , and after triparting the system becomes  $-|111\rangle$ . Therefore,  $|1\rangle|00\rangle \rightarrow \frac{1}{\sqrt{2}}(|000\rangle - |111\rangle) = |-\rangle$ .

10.5 Suppose the first qubit is flipped,  $x'_1 = x_1 \oplus 1$ ; so  $x'_1 \oplus x_2 = 1 \oplus 0 = 1, x'_1 \oplus x_3 = 1 \oplus 0 = 1$ . That is,  $A_1 = 1, B_1 = 1$ .

Suppose the second qubit is flipped,  $x'_2 = x_2 \oplus 1$ ; so  $x_1 \oplus x'_2 = 0 \oplus 1 = 1, x_1 \oplus x_3 = 0$ . That is,  $A_1 = 1, B_1 = 0$ .

Suppose the third qubit is flipped,  $x'_3 = x_3 \oplus 1$ ; so  $x_1 \oplus x_2 = 1, x_1 \oplus x'_3 = 0 \oplus 1 = 1$ . That is,  $A_1 = 1, B_1 = 1$ .

10.6 When there is no error,  $x_1 \oplus x_2 = x_1 \oplus x_3 = 0$

Because there is no error in the second and third group,  $A_2 = B_2 = 0, A_3 = B_3 = 0$ .

When a phase error exists in the second qubit  $x'_2 = -x_2$ . We see  $x_2 - x'_2 = 0 \pmod{2}$ .  $x_1 \oplus x'_2 = 0 - 0 \pmod{2} = 0$ . And because  $x_3$  is unchanged, we still have  $x_1 \oplus x_3 = 0$ .

Therefore,  $A_1 = B_1 = 0$ .

10.7 (1) Because there is no error in the first and second qubit,  $A_1 = B_1 = 0, A_2 = B_2 = 0$ . As we have seen in part 10.6, a phase shift does not vary  $A_i, B_i$ , effect of a  $Y$  error on  $A_3, B_3$  is equivalent to an  $X$  error on them. As found in part 10.5,  $A_3 = B_3 = 1$ .

(2) Applying  $\sigma_x$  to first qubit of the third group, we obtain  $a|++\rangle + b|--\rangle$ . From table 10.3,  $(A_4, B_4) = (0, 1)$ .

(3) Note that  $Y = ZX$ . So,

$$\begin{aligned} & (U_\beta \otimes I \otimes I)(a|000\rangle + b|111\rangle) = \cos \beta(a|000\rangle + b|111\rangle) + i \sin \beta(a|000\rangle - b|011\rangle), \text{ and} \\ & (U_\alpha \otimes I \otimes I)(a|000\rangle + b|111\rangle)(\cos \beta(a|000\rangle + b|111\rangle) + i \sin \beta(a|000\rangle - b|011\rangle)) \\ & = \cos(\alpha + \beta)(a|000\rangle + b|111\rangle) + i \sin(\alpha + \beta)(a|100\rangle + b|011\rangle). \end{aligned}$$

Thus, we have  $(A_4, B_4) = (1, 1)$  after applying the first layer of the Hadamard gates.

10.8 Since  $H[0001110]^t = [3 \pmod{2}, 1, 1]^T = [1, 1, 1]^t$ ,  $H[1101000]^t = [1, 1, 1]^t$ , and

$$H[1100111]^t = [3 \pmod{2}, 3 \pmod{2}, 3 \pmod{2}]^t = [1, 1, 1]^t,$$

we apply a bit flip at the seventh bit in each case to do the correction.

10.9 (a) Since  $d(c_1, c_2)$  is the number of different digits in  $c_1, c_2$ ,  $d(c_1, c_2) \geq 0$ .

(b) Since the number of digits in  $c_2$  differing from  $c_1$  is the same as the number of digits in  $c_1$  differing from  $c_2$ , we see that  $d(c_1, c_2) = d(c_2, c_1)$ .

(c) Let  $c_1 = (x_1, \dots, x_n), c_2 = (y_1, \dots, y_n), c_3 = (z_1, \dots, z_n)$ . Define  $|u - v| = 0$  if  $u = v$  and  $|u - v| = 1$  if  $u \neq v$ . Then  $d(c_1, c_2) = \sum_j |x_j - y_j|$ ,  $d(c_2, c_3) = \sum_j |y_j - z_j|$ . It is clear that for each  $j$ ,  $|x_j - z_j| \leq |x_j - y_j| + |y_j - z_j|$ . So,  $d(c_1, c_3) = \sum_j |x_j - z_j| \leq d(c_1, c_2) + d(c_2, c_3)$ .

10.10 Note that  $d(c_1, c_2) = d(c_1 - c_1, c_2 - c_1)$  because both sides equal the numbers of entries in  $c_1$  differing from  $c_2$ . So,  $\min\{d(c_1, c_2) : c_1 \neq c_2\}$  is the same as  $\min\{d(0, c) : c \neq 0\}$ . One readily see that the 15 nonzero codewords has at least 3 nonzero entries. The result follows.

10.11 (1) Note that  $X^2 = Z^2 = I_2$ . Thus,  $M_i^2 = I$  for  $i = 0, \dots, 4$ . Because  $M_i$  is a tensor product of  $I_2, X, Z$  and all these matrices have eigenvalues in  $\{1, -1\}$ , the eigenvalues of  $M_i$  will be product of numbers in  $\{1, -1\}$  and will lie in the set  $\{1, -1\}$ .

(2) Use the fact that  $XZ = -ZX, X^2 = Z^2 = I$ . We have

$$\begin{aligned} [M_0, M_1] &= Z \otimes Z \otimes XZ \otimes X^2 \otimes ZX - Z \otimes Z \otimes ZX \otimes X^2 \otimes XZ = 0, \\ [M_0, M_2] &= X \otimes Z^2 \otimes X \otimes XZ \otimes ZX - X \otimes Z^2 \otimes X \otimes ZX \otimes XZ = 0, \\ [M_0, M_3] &= X \otimes ZX \otimes XZ \otimes X \otimes Z^2 - X \otimes XZ \otimes ZX \otimes X \otimes Z^2 = 0, \\ [M_1, M_2] &= ZX \otimes Z \otimes Z \otimes XZ \otimes X^2 - XZ \otimes Z \otimes Z \otimes ZX \otimes X^2 = 0, \\ [M_1, M_3] &= ZX \otimes X \otimes Z^2 \otimes X \otimes XZ - XZ \otimes X \otimes Z^2 \otimes X \otimes ZX = 0, \\ [M_2, M_3] &= X^2 \otimes ZX \otimes Z \otimes Z \otimes XZ - X^2 \otimes XZ \otimes Z \otimes Z \otimes ZX = 0. \end{aligned}$$

$$\begin{aligned} (3) M_4 &= M_0 M_1 M_2 M_3 = (X_2 X_3 Z_1 Z_4)(X_3 X_4 Z_2 Z_0)(X_4 X_0 Z_3 Z_1)(X_0 X_1 Z_4 Z_2) \\ &= X_2 X_3^2 Z_1 Z_4 X_4^2 Z_2 Z_0 X_0^2 Z_3 Z_1 X_1 Z_4 Z_2 = X_2 Z_1^2 Z_4^2 Z_2^2 Z_0 Z_3 Z_1 = X_1 X_2 Z_0 Z_3. \end{aligned}$$